

# Weighted LARS for Quantitative Stream Reasoning

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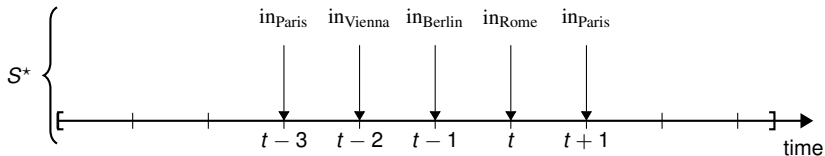
12<sup>th</sup> of September 2020

**FWF**

Der Wissenschaftsfonds.

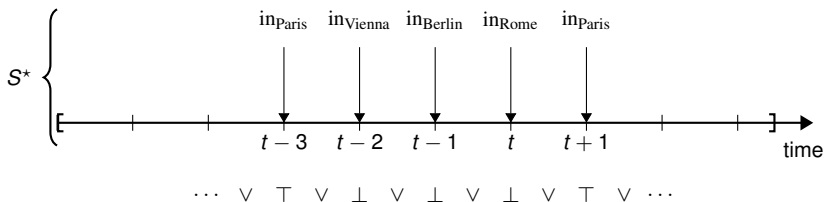
**logics**  LOGICAL METHODS IN  
COMPUTER SCIENCE

# (Qualitative) Stream Reasoning with LARS



- ▶ Does Peter visit Paris?

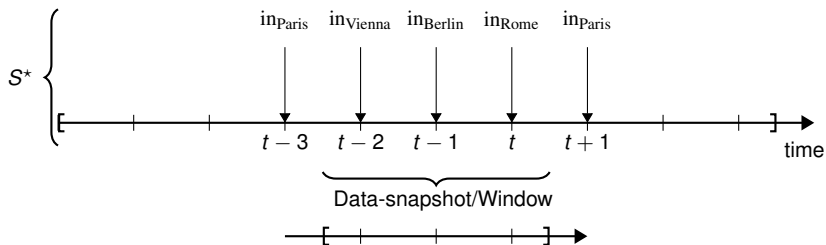
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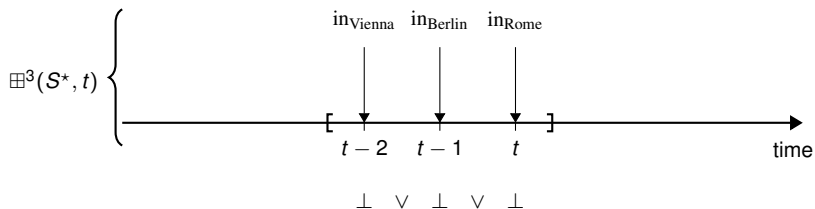


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- ▶ Other quantitative questions (Probabilities, Weighted Model Counting, ...)

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- ▶ Weighted Logic!
- ▶ *Overload* ( $\{\perp, \top\}, \vee, \wedge, \perp, \top$ ) using semirings  $(R, \oplus, \otimes, e_{\oplus}, e_{\otimes})$  and allow semiring values in formulas

# Semiring Semantics

Examples are

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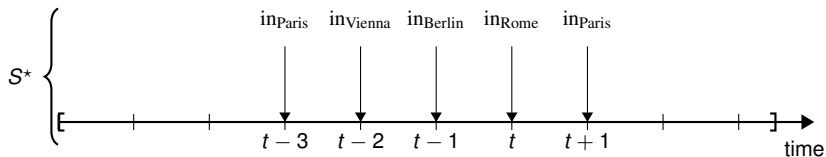
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↪ classical semantics
- ▶  $\mathbb{S} = (\mathbb{S}, +, \cdot, 0, 1)$ , for  $\mathbb{S} \in \{\mathbb{N}, \mathbb{Z}, \mathbb{Q}, \mathbb{R}\}$ , the semiring over the numbers in  $\mathbb{S}$ .  
↪ disjunction is sum, conjunction is multiplication

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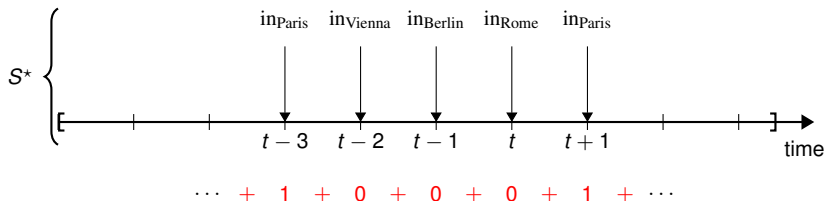
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     $\hookrightarrow$  disjunction is sum, conjunction is multiplication
- ▶  $\mathcal{R}_{\max} = (\mathbb{Q} \cup \{-\infty, \infty\}, \max, +, -\infty, 0)$ , the max-tropical semiring.  
     $\hookrightarrow$  disjunction is maximum, conjunction is sum

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→  $\diamond \text{in}_{\text{Paris}}$  over the natural number semiring  $(\mathbb{N}, +, \cdot, 0, 1)$

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$$\mu(S, t) = \begin{cases} \llbracket \alpha \rrbracket_{\mathcal{R}}(S, S, t) & \text{if } S \text{ is an answer stream of } \Pi \text{ at } t, \\ \mathbf{e}_{\oplus} & \text{otherwise.} \end{cases}$$

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- ▶  $S$  is an answer stream of  $\Pi$  at  $t$  if  $(S, S, t)$  satisfies  $\Pi$  and  $(S, S, t)$  is a minimal model of the reduct  $\Pi^{S,t} = \{\alpha \leftarrow \beta \in \Pi \mid (S, S, t) \text{ satisfies } \beta\}$



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▶ **Weighted Model Counting** by

computing the sum using  $\oplus$   
of  $\mu(S, t)$  for each answer streams  $(S, t)$

## Expressivity Results I

LARS measures enable subsumption of corresponding ASP-extensions

- ▶ Preferential Reasoning, i.e.,
  - Weak Constraints [Buccafurri *et al.*, 2000]
  - (part.) *asprin* [Brewka *et al.*, 2015]
- ▶ Probabilistic Reasoning, i.e.,
  - P-log [Baral *et al.*, 2009]
  - LP<sup>MLN</sup> [Lee and Yang, 2017]
  - ProbLog [De Raedt *et al.*, 2007]
- ▶ Weighted Model Counting, i.e.,
  - aProbLog [Kimmig *et al.*, 2011]

## Expressivity Results II

- ▶ A plain fragment of LARS measures is expressively equivalent to
  - ▶ Weighted Automata (Finite State Machines with weighted transition function)
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- ▶ Gives a rule-based alternative for specification via automata

## Complexity Results

The evaluation of LARS measures

- ▶ is PSPACE-hard for any non-trivial semiring (LARS is already PSPACE-complete)
- ▶ possible in FPSPACE(poly) for under mild restrictions on the semiring and weighted formula

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Preferential Reasoning (over restricted LARS measures)

- ▶ Preference Checking is  $\Pi_2^P$ -complete
- ▶ Brave Preferential Reasoning is  $\Sigma_3^P$ -complete

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# Conclusion & Outlook

- ▶ LARS enables expressive stream reasoning
- ▶ Weighted LARS and LARS measures as a general underlying framework for quantitative stream reasoning
  - ↔ Lift quantitative LP-extensions to the streaming context
- ▶ Restrictions on LARS measures can tame the complexity
- ▶ Next up
  - ▶ Implementation
  - ▶ Application in object detection, traffic regulation





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# Weighted LARS

- ▶ Evaluate formulas w.r.t.  $(S^*, S, t)$
- ▶  $S^*$  starting stream,  $S$  current stream,  $t$  time

Formula	LARS	Weighted LARS
constant	$\perp, \top$	semiring value $k$
$p$	true, false	one, zero
$\neg\alpha$	true $\leftrightarrow$ false	zero $\rightarrow$ one, rest $\rightarrow$ zero
$\alpha \wedge \beta$	$\alpha$ and $\beta$	$\alpha$ times $\beta$
$\alpha \vee \beta$	$\alpha$ or $\beta$	$\alpha$ plus $\beta$
$\Box\alpha$	for all $t: \alpha$	product of $\alpha$ over $t$
$\Diamond\alpha$	exists $t: \alpha$	sum of $\alpha$ over $t$
$\textcircled{t}\alpha$	$(S^*, S, t)$ changes to $(S^*, S, t')$	
$\boxplus^w\alpha$	$(S^*, S, t)$ changes to $(S^*, w(S, t), t)$	
$\triangleright\alpha$	$(S^*, S, t)$ changes to $(S^*, S^*, t)$	