

Reasoning on Contextual Hierarchies via ASP with Algebraic Measures

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January 2021 - May 2021

Context Representation

- Context representation is a well-known area of study in KR
- Semantic Web: need to interpret datasets in correct context
- Different proposals for DL-based representation of contexts

Contextualized Knowledge Repositories (CKR)

DL-based framework for reasoning with contextual knowledge in the semantic web [2]

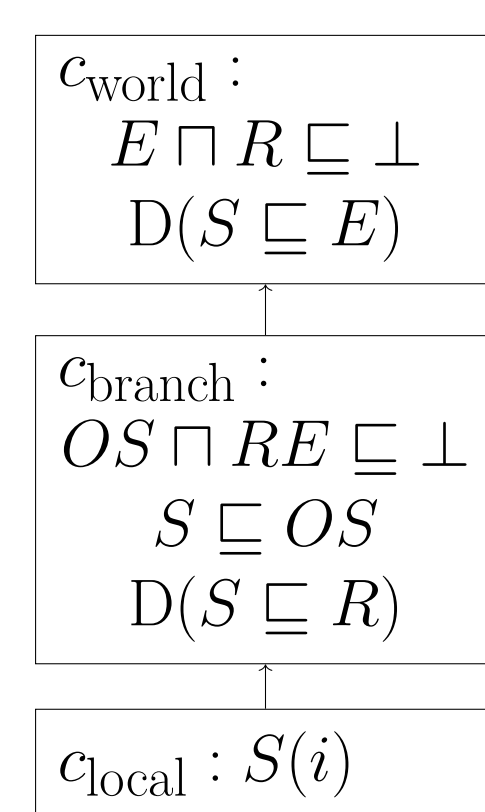
A CKR $\mathfrak{K} = \langle \mathcal{C}, K_{\mathbf{N}} \rangle$ consists of

- Global Knowledge $\mathcal{C} = \langle \mathbf{N}, \succ \rangle$:
The relation \succ between the different contexts $c \in \mathbf{N}$
- Contextual Knowledge $K_{\mathbf{N}} = (K_c)_{c \in \mathbf{N}}$:
The *additional* axioms K_c in context $c \in \mathbf{N}$

In local contexts we have, apart from usual DL rules

- Defeasible Axioms $D(C \sqsubseteq D)$:
Can be overridden in more specific contexts
- Eval-Expressions $eval(c, D)$:
Can reference the state of D in another context c

CKR Example



- Electronics, Robotics, Supervisor
- OnSite, REMote
- At c_{local} we have $S(i)$, $OS(i)$ and $R(i)$
- $S \sqsubseteq OS$ is actually defeasible w.r.t. time!
- $D(S \sqsubseteq R)$ actually only holds since 2020!

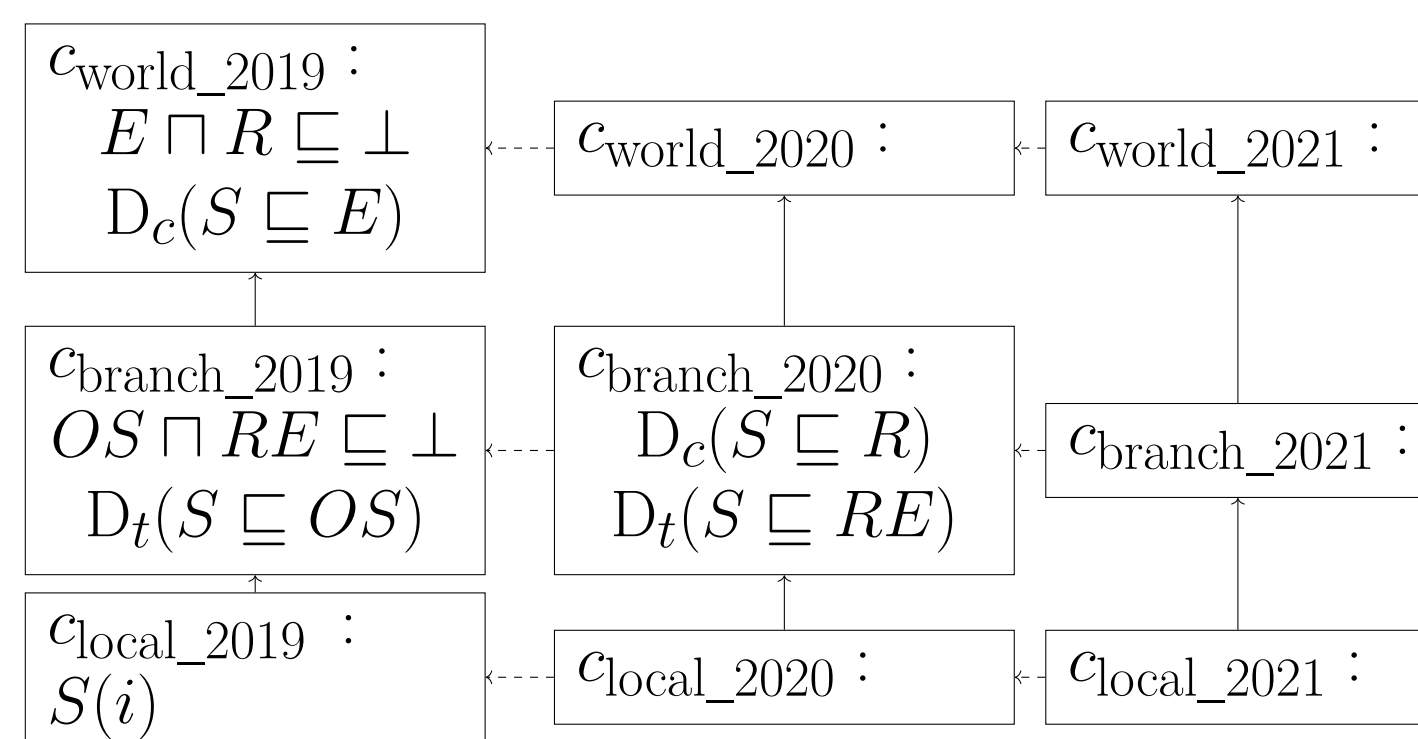
Multi-Relational CKR

- We need different relations between contexts!
- Axioms may be defeasible *with respect to a relation!*

- Global Knowledge $\mathcal{C} = \langle \mathbf{N}, \succ_1, \dots, \succ_m \rangle$:
The relations \succ_i between the different contexts $c \in \mathbf{N}$
- Contextual Knowledge $K_{\mathbf{N}} = (K_c)_{c \in \mathbf{N}}$:
The *additional* axioms K_c in context $c \in \mathbf{N}$

- Defeasible Axioms $D_i(C \sqsubseteq D)$:
Can be overridden in more specific contexts *w.r.t.* \succ_i

MR-CKR Example



- ... for \succ_c , the coverage relation
- ... for \succ_t , the time relation
- At c_{local_2019} we have $S(i)$, $OS(i)$ and $E(i)$
- At c_{local_2020} and c_{local_2021} we have $S(i)$, $RE(i)$ and $R(i)$

Semantics

- Clashing assumption $\langle \alpha, e \rangle$: instance e is an exception of $D_r(\alpha)$
- CAS-interpretation $\mathcal{I}_{CAS} = \langle \mathcal{I}, \chi_1, \dots, \chi_m \rangle$:
 - $\mathcal{I}(c)$: interpretation of context c
 - $\chi_i(c)$: set of clashing assumptions of context c w.r.t. relation \succ_i

(Justified) CAS-model $\mathcal{I}_{CAS} \models \mathfrak{K}$

\mathcal{I}_{CAS} is a CAS-model for \mathfrak{K} if:

- $\mathcal{I}(c) \models K_c$, if $c' \preceq c$
- for every $D_i(\alpha) \in K_c$ and $c' \preceq_i c$, $\mathcal{I}(c') \models \alpha$
- for every $D_i(\alpha) \in K_c$ and $c' \preceq_i c$, if $\langle \alpha, e \rangle \notin \chi_i(c')$, then $\mathcal{I}(c') \models \alpha(e)$

\mathcal{I}_{CAS} is justified if each clashing assumption $\langle \alpha, e \rangle \in \chi(c)$ is justified by some clashing set S such that

- $\mathcal{I}(c) \models S$
- $S \cup \{\alpha(e)\}$ is unsatisfiable

Which justified CAS-Models are preferred?

(LP) Locally, we prefer those that satisfy more specific defeasible axioms: $\chi_i(c) > \chi'_i(c)$, if

- for every $\eta = \langle \alpha, e \rangle \in \chi_i(c) \setminus \chi'_i(c)$ with $D_i(\alpha)$ at context $c_1 \succeq_i c_{1b} \succ_i c$,
- there exists $\eta' = \langle \alpha', e' \rangle \in \chi'_i(c) \setminus \chi_i(c)$ with $D_i(\alpha')$ at context $c_2 \succeq_i c_{2b} \succ_i c$
- such that $c_{1b} \succ_i c_{2b}$

(RP) On the relation level, we prefer those that have an improvement locally and no change for the worse otherwise: $\chi_i > \chi'_i$, if

- there exists $c \in \mathbf{N}$ s.t. $\chi_i(c) > \chi'_i(c)$ and not $\chi'_i(c) > \chi_i(c)$
- for no context $c' \neq c \in \mathbf{N}$ it holds that $\chi_i(c') < \chi'_i(c')$ and not $\chi'_i(c') < \chi_i(c')$.

(GP) Globally, we prefer those that are preferred on the relation of the smallest index:

- $\langle \mathcal{I}, \chi_1, \dots, \chi_m \rangle > \langle \mathcal{I}', \chi'_1, \dots, \chi'_m \rangle$, if
- there exists $i \in \{1, \dots, m\}$ such that $\chi_i > \chi'_i$
- for all $j < i \in \{1, \dots, m\}$ it holds that $\chi_j \not< \chi'_j$

Reasoning with MR-CKR

How can we reason with MR-CKR?

- Restrict the DL language to *SROIQ-RL* and obtain least models via translation to ASP [1]

But how can we get only the preferred models?

- Preferences + ASP \rightarrow asprin [3]
- asprin only allows strict partial orders but we can have cyclic preference relations $\not\prec$

We investigated two options:

- Restriction to *eval-disconnected* MR CKR: avoid cycles
- Use algebraic measures

eval-Disconnectedness

Idea:

- If \mathfrak{K} is *eval-free*, i.e., there are no *eval-expressions* at all, the interpretations $\mathcal{I}(c)$ and $\mathcal{I}(c')$ for $c \neq c'$ are independent
- Then any interpretation $(\mathcal{I}(c))_{c \in \mathbf{N}}$, where $\mathcal{I}(c)$ is locally preferred, is also globally preferred

We want at least some *eval-expressions* though!

Introduce *eval-Disconnectedness*

- Generalizes this idea
- Is a syntactic condition that can be checked easily
- Prevents dependence of the satisfaction of a default at context c on the satisfaction of another default at context c'

asprin Encoding

Local Preference (LP)

```
#preference(LP(c,i),poset){
  -ovr(alpha,X,c,i) >> ovr(alpha,X,c,i);
  -ovr(alpha2,Y,c,i) >> -ovr(alpha1,X,c,i);
  for c1 >=i c1b >_i c and
  c2 >=i c2b >_i c and c1b >_i c2b and
  Di(alpha) in Kc. }
```

Relation-global Preference (RP)

```
#preference(RP(i),pareto){
  **LP(C,i) : context(C) }.
```

GlobalPreference (GP)

```
#preference(GP,lexico){
  W:**RP(I) : rel_w(I,W) }.
```

Correctness

Putting things together:

- $PK(\mathfrak{K})$ is the answer set program that encodes the MR-CKR \mathfrak{K}
- P_{pref} is the preference encoding in asprin including #optimize(GP).

Theorem

Let \mathfrak{K} be a multi-relational CKR that is *eval-disconnected* and in *SROIQ-RLD normal form*. Then under the unique name assumption,

- for every α and c such that $O(\alpha, c)$ is defined, $\mathfrak{K} \models c : \alpha$ iff $PK(\mathfrak{K}) \cup P_{pref} \models O(\alpha, c)$;
- for every *BCQ* $Q = \exists y. \gamma(y)$ on \mathfrak{K} , $\mathfrak{K} \models Q$ iff $PK(\mathfrak{K}) \cup P_{pref} \models O(Q)$.

We can use $PK(\mathfrak{K}) \cup P_{pref}$ to reduce reasoning tasks to ASP+asprin!

\hookrightarrow implemented in the publicly available tool **CK-Rew**!

Weighted Logic [4]

- Use weighted formulas α over a semiring $(R, \oplus, \otimes, e_{\oplus}, e_{\otimes})$ of the form

$$\alpha ::= k \mid p \mid \neg p \mid \alpha + \alpha \mid \alpha * \alpha,$$

where $k \in R$ and p is a propositional variable

- Allows calculations over a semiring depending on the truth of propositional variables or formulas
- Example: $2 * \text{candy} + 3 * \text{pasta}$ over semiring $(\mathbf{N}, +, \cdot, 0, 1)$
 \hookrightarrow if we buy candy and pasta we pay 5

Algebraic Measures

- An Algebraic Measure μ is defined by a triple $\langle \Pi, \alpha, \mathcal{R} \rangle$, where
 - Π is an ASP program
 - α is a weighted LARS formula over \mathcal{R}
 - \mathcal{R} is a semiring

- The weight of an answer set $S \in \mathcal{AS}(\Pi)$ is

$$\mu(S) = \llbracket \alpha \rrbracket_{\mathcal{R}}(S).$$

- Intuitively, algebraic measures allow us to associate a weight with an answer set \hookrightarrow many possibilities what to do with this weight!

- The overall weight of μ is defined as

$$\mu(\Pi) = \bigoplus_{S \in \mathcal{AS}(\Pi)} \mu(S).$$

- An answer set $S \in \mathcal{AS}(\Pi)$ is preferred w.r.t. a measure $\mu = \langle \Pi, \alpha, \mathcal{R} \rangle$ and a relation $>$ on R if no $S' \in \mathcal{AS}(\Pi)$ exists such that $\mu(S') > \mu(S)$

Preference Encoding

We can use this as follows

- The powerset semiring $\mathcal{P}(CA)$ over the set of possible clashing assumption can do “bookkeeping”
- Define a weighted formula that checks for clashing assumptions
 $\alpha = \sum_{\langle \phi, e, c, i \rangle \in CA} \text{ovr}(\phi, e, c, i) * \{ \langle \phi, e, c, i \rangle \}$.
- Take the relation $>_{opt}$ on the semiring values $S \subseteq CA$ that correctly captures the preference on the justified models.
- $S \in \mathcal{AS}(PK(\mathfrak{K}))$ is preferred iff it corresponds to a least preferred CAS model $\langle \mathcal{I}, \bar{\chi} \rangle$ of \mathfrak{K}

Conclusions

- Multi-Relational CKR allows us to properly capture differences in contexts w.r.t. different dimensions
- CKRw: Implementation using an encoding in ASP + asprin for *eval-disconnected* MR-CKR
- Algebraic Measures as a general approach many quantitative problems
 \hookrightarrow probabilistic reasoning
 \hookrightarrow preferential reasoning
 \hookrightarrow parameter learning
 \hookrightarrow and more

References

- Loris Bozzato, Thomas Eiter, and Luciano Serafini. Enhancing context knowledge repositories with justifiable exceptions. *Artif. Intell.*, 257:72–126, 2018.
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1 Paper and Implementation

Paper: “Reasoning on Multi-Relational Contextual Hierarchies via Answer Set Programming with Algebraic Measures,” ICLP 2021.
Prototype & Tech. Report: available at <https://www.ai4europe.eu/research/ai-catalog/ckrew-ckr-datalog-rewriter>